

Patent Application of
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For

TITLE: MULTI-MODE TOKEN

CROSS-REFERENCE TO RELATED APPLICATIONS: Not Applicable

FEDERALLY SPONSORED RESEARCH: Not Applicable

SEQUENCE LISTING OR PROGRAM: Not Applicable

BACKGROUND OF THE INVENTION – FIELD OF INVENTION

[001] This invention relates to authentication token devices such as Integrated Circuit (IC) cards, Universal Serial Bus (USB) keys, Biometric Authentication devices, Remote Access Key Tokens, or Cellular Phones.

BACKGROUND OF THE INVENTION

[002] This invention provides multiple authentication modes for authentication devices such as IC Cards, USB keys, Biometric Authentication devices, Remote Access Key Tokens, or Cellular Phones depending on the situation.

[003] Current conventional authentication token devices have at most two modes. One mode does not require a Personal Identification Number (PIN), Password or Biometrics such as a fingerprint or iris.

The other mode requires a PIN, Password or Biometrics at every event.

[004] For many applications, it is inconvenient to only have these two modes available.

[005] For example,

many companies would like employees to use a single IC card to pass through the physical gates, access computers, access the company’s computer servers, access high security rooms, make small payments at the company cafeteria or make authorizations using digital signatures.

Each of these applications requires a different level of security and increasing security decreases user convenience.

SUMMARY

[006] The object of this invention, Multi-Mode Token, is to provide at least one mode that allows for a certain period of time or number of operations to pass before requiring a user to logon in order to balance security and convenience, solving the aforementioned problem **[005]**.

BRIEF DESCRIPTION OF THE PROCESS FLOW AND TABLES – FIGURES

[007] FIG 1 is a table of the notation used in this application.

FIG 2 is the formulae of this application.

FIG 3 is a block diagram of the system of this invention.

FIG 4 is an example of the modes of this invention, Multi-Mode Token.

FIG 5 is a table of Register and Counter values in a Multi-Mode Token.

FIG 6 is a table of the basic user data stored on a Multi-Mode Token.

FIG 7 is an example of multi-mode settings.

FIG 8A is the initialization flow of a Multi-Mode Token.

FIG 8B is the initialization flow of a Multi-Mode Token continued.

FIG 9 is the flow of Multi-Mode Token logon.

FIG 10 is the flow of mode 1 operation.

FIG 11 is the flow of mode 2, authentication.

FIG 12 is the flow of mode 2, decryption.

FIG 13 is the flow of mode 2, payment.

FIG 14 is the flow of mode 3 payment/authorization

FIG 15 is the flow of mode 4 payment/authorization.

DETAILED DESCRIPTION

[008] Notation

FIG 1 outlines the notation used in this application. Bold uppercase letters represent entity names such as **A**, **B**, or items such as an ID# **N**, keys **C**, **D**, **E** or times **TP**, **TL**. Lowercase bold suffixes signify a relation to the relative entity.

[009] Formulae

FIG 2 lists all formulas used in this application.

[010] System Block Diagram

FIG 3 provides an embodiment of the system of this invention in the form of a block diagram.

This system includes at least one authenticator module **A_j** (312) at a local system terminal **J** (310) and at least one authentication token **B_i** (322) owned by user **I** (320) and a certificate module **A_o** (302) at the system Authority **O** (300). **A_o** and **A_j** have a token interface.

User **I** can input his/her unique feature **F_i** into the token **B_i** via an input device (330). The unique feature **F_i** might be **F_{1i}**: the user's PIN or password, or **F_{2i}**: a biometric feature such as a fingerprint, hand shape, iris, face or voice.

In most cases, the input device (330) is provided by Terminal **J**, but some tokens include an integrated biometric sensor or keypad for PIN input.

[011] Conventional System

Conventional token authentication systems have only two modes.

The first mode is used for relatively low security access and does not require users to log on to the token.

The second mode is used for higher security access and requires users to log on to the token using a unique feature **F_i**, such as a PIN, password or biometrics, at every session.

[012] Examples of Problems with the Conventional System

- a) Employees must possess a token in order to pass through the company entrance or parking gate.

It is too risky for employees to be allowed to pass through the gates using only a token because the token might be lost or stolen.

In this case, it is necessary to have at least a minimal relationship between users and their tokens.

But if all employees are required to log on at a physical gate, the gate would be jammed due to the time it takes for employees to log on.

In this case, one day is too short because every Monday morning the logon status of most tokens would be expired.

Therefore, it is desired that one logon be valid for one week at the company gates.

- b) Employees use tokens to pay for meals at the company cafeteria.

Requiring employees to log on at each cafeteria counter is tedious and would cause lines to congest.

On the other hand, it is too risky to make logons valid for a week.

Therefore, it is desired that one logon be valid for one day in the company cafeteria.

[013] Multi-Mode Token

In order to solve problems like the above [012], this invention provides a Multi-Mode Token.

It is ideal to use one token for multiple applications which all require different access security levels.

FIG 4 illustrates one example of mode usage for multiple applications which all require different access security levels.

FIG 5 shows the necessary register and counters that are used in a Multi-Mode Token.

The essential concept of this invention is to provide those registers in a token and to use them if required by the security level of the application.

[014] FIG 6 shows a table of necessary data contained in a Multi-Mode Token.

Some of this data may be omitted depending on the system in which the token is used. The following sections describe how a token obtains and uses this data. The common key **C**, private decryption key **D**, private signing key **S**, hash value of the password **H1**, and the feature vector of biometrics **H2** are kept secret and are used only in certain modes.

[015] Mode Settings

FIG 7 provides a more detailed example of the mode settings seen in FIG 4.

Users are not required to log on to the token for mode 0 operation, in which the application obtains the non-secret information seen in FIG 6.

In the mode 1 operation for the relatively low security application, the common secret key **C** is used and user logon is required, which is valid for one week or for ten operations.

In the mode 2 operation for the second-level security application, the private decryption key **D** is used and user logon is required, which is valid for one day or for five operations.

In modes 1 and 2, the relationship between the Expiration Date Condition

$$\text{formula (203)} \quad TP - TL \leq TM$$

and the Access Times Condition

$$\text{formula (204)} \quad G > 0$$

can be either the “OR” condition or the “AND” condition.

In the mode 3 operation for the high-level security application, user logon is required for every operation using **S** since the maximum of the **G3** and **G4** counters is 1, which is reduced by 1 at each operation using **D** or **S**.

In the mode 4 operation for the highest-level security application, two different types of logon, such as a password and fingerprint, are required for every operation using **S**.

G1, G2 and G3 are reset to **G max** at every logon using **F1** or **F2**.

G4 is reset to 1 at every logon using **F2** after **G3** is reset by logging on with **F1**.

[016] Initialization

FIG 8A shows the initialization flow of a token.

Initialization is performed by the certificate module **Ao** at the System Authority or System Administrator **O**.

Flow (800) – In preparation, **Ao** generates a System Master Common Key **Co** and a signing key pair **So, Vo**.

Flow (801) – User Token **Bi** obtains a User Name, a hash value of User Password **H1i**, feature vectors of User Biometrics **H2i** from User **I** by some secure means, and obtains ID# **Ni** and the initial **Ci** from **Ao** by some secure means.

Flow (804) – A session is setup between **Ao** and **Bi**.

Depending on the situation, the old **Ci** may be used as the session key for all secure communication within the initial session.

Flow (806) – **Bi** requests a new **Ci** from **Ao**, sending **Ni** in order to validate the token.

Flow (807) – **Ao** derives a new **Ci** using

$$\text{formula (205)} \quad \mathbf{Ci} = \mathbf{Co} \{ \mathbf{Ni} + \mathbf{TC} \}.$$

The expiration date **TC** of the key **Ci** is embedded into **Ci** itself so that **Ci** cannot be used after that date.

Flows (808) and (809) – **Ci** and **TC** are returned to and stored on **Bi**. **Ci** is used in mode 1.

Flow (816) – In the next initialization step, **Bi** sends **Ni** to **Ao** and requests an asymmetric key pair **Di, Ei** and its certificate **LEi**.

Flow (817) – **Ao** generates the asymmetric key pair **Di, Ei** and its certificate **LEi** using

$$\text{formula (207)} \quad \mathbf{LEi} = \mathbf{So} \{ \mathbf{Ni}, \mathbf{Ei}, \mathbf{TE} \}.$$

As formula (207) shows, **LEi** is a certificate issued by **Ao** to authorize and validate **Ni** and **Ei** until the expiration date **TE**.

Flow (818) – **Ao** returns **Ei**, **LEi** and **TE** to **Bi**. **Ao** also returns **Di** to **Bi** in secure manner such as wrapping it with the session key. **Di** is also escrowed by **Ao**.

Flow (819) – **Bi** stores **Ei**, **LEi**, **TE** and **Di** and uses them for session key exchange, decryption of file encryption keys, authentication at physical gates, or micro payments.

Flow (823) – **Bi** itself generates a key pair **Si**, **Vi**.

Si is a signing key and **Vi** is a verification key.

Flow (826) – **Bi** sends **Ni** and **Vi** to **Ao** and requests the certificate **LVi**.

Flow (827) – **Ao** calculates the certificate **LVi** using

$$\text{formula (209)} \quad \mathbf{LVi} = \mathbf{So} \{ \mathbf{Ni}, \mathbf{Vi}, \mathbf{TV} \}$$

As formula (209) shows, **LVi** is a certificate issued by **Ao** to authorize and validate **Ni** and **Vi** until the expiration date **TV**.

Flows (828) and (829) – **LVi** and **TV** are returned to and stored on **Bi**. **Si** is used for important signing authorized by User **I**.

TM, **G** max and other Mode settings are configured by either User **I** or System Administrator **O**, depending on the system.

[017] Token Logon

Flow (902) – When **Bi** requires user logon, it jumps from the original flow point to Flow (902).

Flows (904), (906) and (907) – After setting up a session, **Bi** requests that the Authentication Module or Application Module **Aj** of the Terminal **J** indicate “Logon (**F1/F2**)”.

The logon is performed using password **F1**, biometrics **F2** or both depending on the required mode.

Flow (910) – According to the indication on the display of Terminal **J**, User **I** inputs his/her unique feature **Fi** via the input device (330).

Flow (911) – **Bi** verifies the user **I**’s Unique Feature **F1i** or **F2i** using the stored User Authentication Reference **H1i** (Hash Value of Password) or **H2i** (Feature Vector of Biometrics).

If **Fi** is accepted,

Flow (912), (914) – **Bi** sends a “Request Date” command to **Aj** and obtains the Present Date **TP**. This is a unique command not found in the conventional token command set.

Flow (915) – **Bi** resets all registers and counters as shown below:

Log on	Time	Register	TL	=	TP
Mode	1	Counter	G1	=	G1 max = 10
Mode	2	Counter	G2	=	G2 max = 5
Mode	3	Counter	G3	=	G3 max = 1
Mode	4	Counter	G4	=	G4 max = 1

The Logon Time Register is used to check the expiration date.

Mode Counters are used to count the number of operations. In modes 1 and 2, the relationship between the Expiration Period and the Mode Counter is either the “OR” condition or the “AND” condition, depending on the system.

[018] Mode 1

As described in [015] and FIG 7, Mode 1 is used for relatively low security applications such as authorizing physical access through the company entrance or parking gate.

The common secret key **Ci** is used for encryption or signing in Mode 1.

FIG 10 shows an example of the main flow of Mode 1 operation.

Flows (1003), (1004) and (1006) – When User Token **Bi** needs to be authenticated, it sets up a session with Authentication Module **Aj** and sends its ID# **Ni** to **Aj**, requesting authentication.

Flow (1007) – **Aj** generates a challenge message **Qi** using

$$\text{formula (211)} \quad \mathbf{Q_i = NR + TP}$$

which comprises a random number **NR** and the present time **TP**.

Flow (1008) – **Aj** sends **Qi** to **Bi** and requests that **Bi** sign **Qi** using the common secret key **Ci**.

Flow (1011) – **Bi** checks to see if the present time is before the expiration date using

$$\text{formula (203)} \quad \mathbf{TP - TL \leq TM1}$$

where

TP : Present Date

TL : The date of the last logon

TM1 : Mode 1 expiration period

Flow (1012) – **Bi** also checks to see if the Mode 1 Counter **G1** has been exceeded using

$$\text{formula (204)} \quad \mathbf{G1 > 0}$$

Flow (1014) – If $\mathbf{TP - TL \leq TM1}$ and $\mathbf{G1 > 0}$, then **Bi** calculates the response message **Ri** using

$$\text{formula (213)} \quad \mathbf{R_i = C_i \{Q_i\}.$$

Flow (1015) – **Bi** reduces the **G1** Counter value by 1.

Flow (1016) – **Bi** returns the response message **Ri** and the expiration date **TC** of **Ci** to **Aj**.

Flow (1017) – **Aj** derives the secret common key **Ci** of **Bi** using

$$\text{formula (205)} \quad \mathbf{C_i = C_o \{N_i + TC\}}$$

and verifies **Qi** with **Ci** using

$$\text{formula (214)} \quad \mathbf{C_i \{R_i\} \Rightarrow Q_i.$$

Flow (1018) – **Aj** sends the result to **Bi**: Accepted or Denied depending on the result of Flow (1017).

In this example flow, the relationship between **TM1** and **G1** is the “AND” condition.

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However, the relationship between them can be the “OR” condition, or just one of them can be used, depending on the system requirement.

[019] Mode 2, Authentication

Mode 2 is used for mid-level security applications such as the decryption of a session key, decryption of a file encryption key, the signing of a challenge response of authentication, or the signing of micro payments.

The private decryption key **Di** is used as a cryptographic key in Mode 2.

FIG 11 shows an example of the main flow of a Mode 2 authentication operation.

Flows (1103), (1104), (1106) and (1107) – Up to Flow (1107), the flow is the same as in Mode 1.

Flow (1108) – **Aj** requests that **Bi** sign the challenge message **Qi** using the private decryption key **Di** (instead of **Ci** as in Mode 1).

Flows (1111), (1112), (1114) and (1115) – These flows are the same as in Mode 1 except

TM2 is used instead of **TM1**,

G2 is used instead of **G1**, and

Di is used instead of **Ci**.

Flow (1116) – **Bi** returns the response message **Ri** and the certificate **LEi** of the public encryption key **Ei** (instead of **TC** as in Mode 1).

Flow (1117) – **Aj** verifies **LEi** using

$$\text{formula (208)} \quad \mathbf{Vo \{LEi\} \Rightarrow Ni, Ei, TE}$$

and verifies **Qi** using **Ei** and

$$\text{formula (216)} \quad \mathbf{Ei \{Ri\} \Rightarrow Qi}$$

Flow (1118) – **Aj** sends the result to **Bi**: Accepted or Denied depending on the result of Flow (1117).

[020] Mode 2, Decryption

This operation is used to decrypt a session key or unwrap a file encryption key.

Flow (1202) – An Application Module or Authentication Module **Aj** has **Pi**, the ciphertext of the session key or file key **K** encrypted by **Ei** using

$$\text{formula (217)} \quad \mathbf{Pi} = \mathbf{Ei} \{ \mathbf{K} \}.$$

Flow (1208) – **Aj** sends **Pi** and the present time **TP** to **Bi** and requests that **Bi** decrypt **Pi** with **Di**.

Flows (1211), (1212), (1214) and (1215) – These flows are the same as in FIG 11 except

Pi is used instead of **Qi**, and

Mi is used instead of **Ri**.

Flow (1216) – **Bi** returns the plaintext **Mi** to **Aj**.

In this case **Mi** is the key **K**.

[021] Mode 2, Payment

This operation is used to authorize a micro payment (e.g. less than \$10).

Flow (1302) – **Aj** has a payment message (or its hash value) **Mi** which must be authorized by User **I**.

Flow (1307) – **Aj** generates a challenge message **Qi** using

$$\text{formula (212)} \quad \mathbf{Qi} = \mathbf{Mi} + \mathbf{TP}$$

These flows are the same as in FIG 11 except that the payment message **Mi** is used instead of a random number **NR** as in Flow (1107).

[022] Mode 3 Payment/Authorization

Mode 3 is used for high security applications such as signing a payment or making an authorization which cannot be denied later.

Flow (1402) – **Aj** has a message (or its hash value) **Mi** which must be authorized by User **I**.

Flow (1408) – **Aj** sends **Mi** to **Bi** and requests that **Bi** sign on **Mi** with the signing key **Si**.

Flow (1409) – **Bi** checks the bit length of **Mi** which should be less than or equal to 64 bits.

If **Mi** is greater than 64 bits, then this operation should be performed in Mode 4.

Flow (1410) – **Bi** jumps to flow (902) in order to initiate user logon.

Flow (1412) – **Bi** checks the value of the Mode 3 Counter **G3**. This should be **G3** max (= 1) after Flows (902) through (915).

Flow (1414) – **Mi** is signed with **Si** using

$$\text{formula (219)} \quad \mathbf{Ui} = \mathbf{Si} \{ \mathbf{Mi} \}.$$

Flow (1415) – the value of the **G3** counter is reduced by 1 so that **G3** becomes 0 in this case.

Flow (1416) – **Bi** returns **Ui** and the certificate **LVi** to **Aj**.

Flow (1417) – **Aj** verifies **LVi** and **Ui** using

$$\text{formula (210)} \quad \mathbf{Vo} \{ \mathbf{LVi} \} \Rightarrow \mathbf{Ni}, \mathbf{Vi}, \mathbf{TV}$$

and

$$\text{formula (220)} \quad \mathbf{Vi} \{ \mathbf{Ui} \} \Rightarrow \mathbf{Mi}$$

Flow (1418) – **Aj** sends the result to **Bi**: Accepted or Denied depending on the result of Flow (1417).

[023] Mode 4 Payment/Authorization

Mode 4 is used for the applications that require the highest levels of security such as authorization of a large payment or of an important document such as a contract. The flow of FIG 15 is as same as in FIG 14. The only differences are that two logons by **F1** and **F2** are required instead of one as in Flow (1510), **G4** is used instead of **G3** as in Flow (1512), and both the **G3** and **G4** counters are cleared.

[024] Administrator Mode

In addition to the above user modes available in a Multi-Mode Token, an administrator mode can be added, in which only an administrator or the Authenticator **Ao** at System Authority can initiate the mode in order to modify system properties such as the value of **G** max, the Expiration period, etc.

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The User Token **Bi** can authenticate the legitimacy of **Ao** or **Aj** via the standard challenge/response procedure using **Vo**.